

FIG. 1A

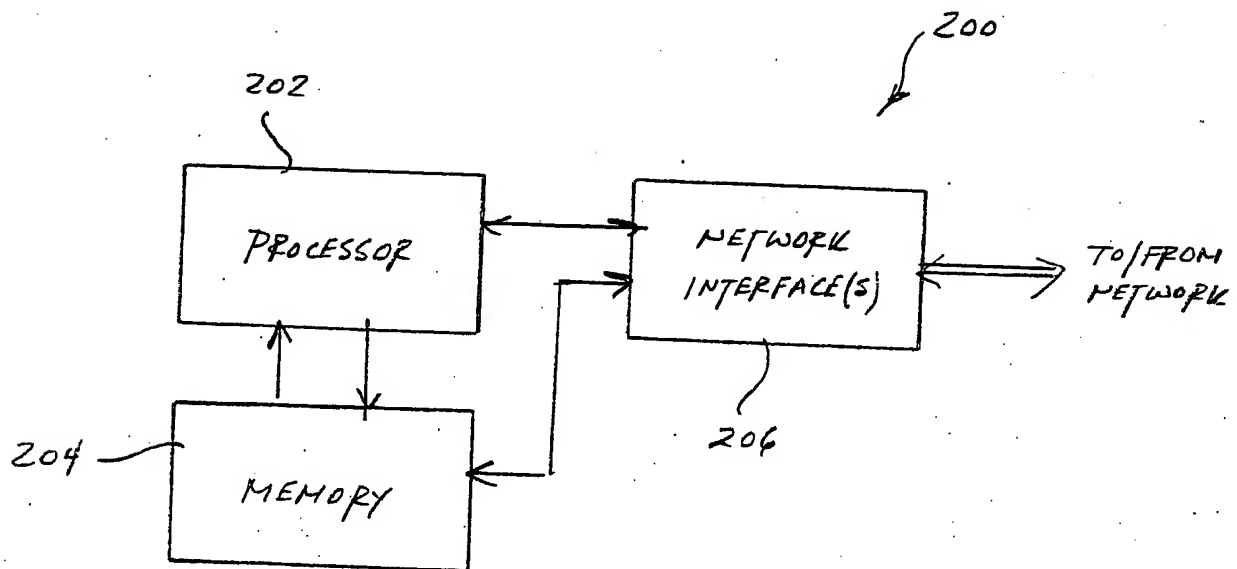


FIG. 1B

222		224	
Paths		Performance Indicator	
Link 1		Clean	
Link 2		Problem	
...		...	
Link n		Clean	

220

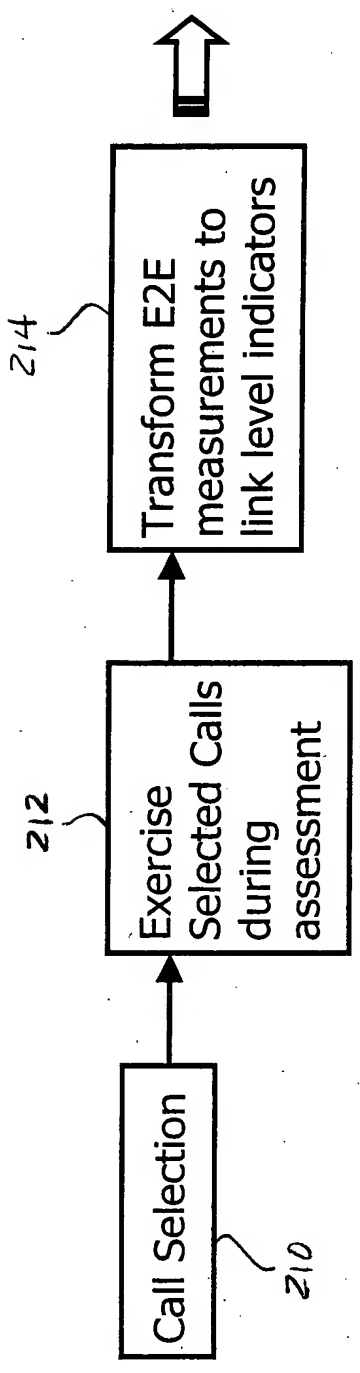


FIG. 2A

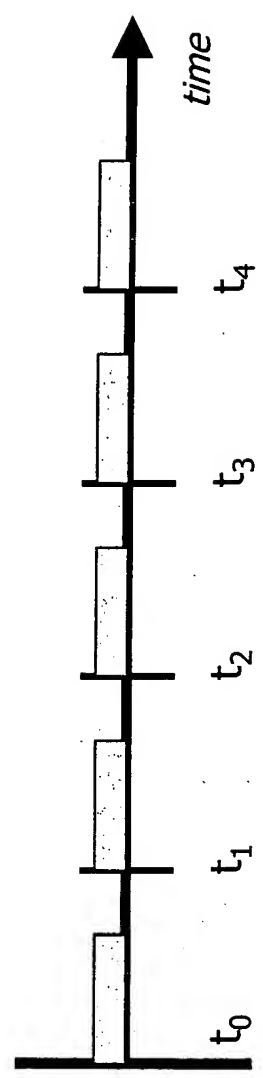


FIG. 2B

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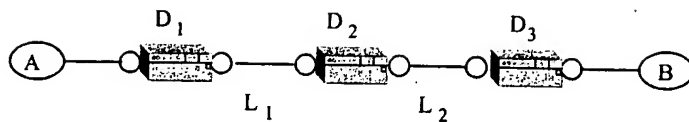


FIG. 3

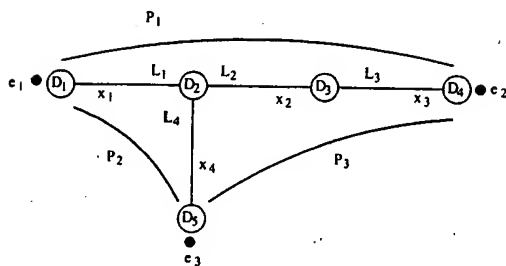


FIG. 4

	L ₁	L ₂	L ₃	L ₄
P ₁	1	1	1	0
P ₂	1	0	0	1

Flow matrix 1

	L ₁	L ₂	L ₃	L ₄
P ₁	1	1	1	0
P ₂	1	0	0	1
P ₃	0	1	1	1

Flow matrix2

FIG. 5

Equations with Flow matrix 1

$$\begin{aligned} x_1 + x_2 + x_3 &= y_1 \\ x_1 + x_4 &= y_2 \end{aligned}$$

$$\begin{bmatrix} 1 & 1 & 1 & 0 \\ 1 & 0 & 0 & 1 \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \\ x_3 \\ x_4 \end{bmatrix} = \begin{bmatrix} y_1 \\ y_2 \end{bmatrix}$$

Equations with Flow matrix 2

$$\begin{aligned} x_1 + x_2 + x_3 &= y_1 \\ x_1 + x_4 &= y_2 \\ x_2 + x_3 + x_4 &= y_3 \end{aligned}$$

$$\begin{bmatrix} 1 & 1 & 1 & 0 \\ 1 & 0 & 0 & 1 \\ 0 & 1 & 1 & 1 \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \\ x_3 \\ x_4 \end{bmatrix} = \begin{bmatrix} y_1 \\ y_2 \\ y_3 \end{bmatrix}$$

FIG. 6

Generate_Pipes($G = (D, L)$): Network Topology Graph, E : Set of Leaves)

I : Set of pipes in G wrt E

$I \leftarrow \emptyset$

Compute P for G wrt E

Let M be the complete flow matrix for G and P

// Group links with the same column vector into disjoint sets

Let k be the number of distinct column vectors in M

Form a set $S = \{S_0, S_1, \dots, S_k\}$ where :

each $S_i, 0 \leq i \leq k$ contains links in L with the i^{th} distinct column vector in M

// Ensure that links in each element of S form a path in G

for $i=1$ to $|S|$

if links in S_i are consecutive and form a path

then merge S_i into path $p, I \leftarrow I \cup \{p\}$

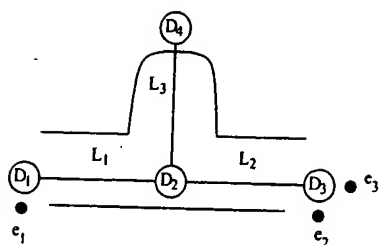
else $I \leftarrow I \cup S_i$

return I

// add the path formed by the links in S_i as a pipe
// add each link as a pipe by itself

FIG. 7

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	L_1	L_2	L_3
$e_1 - e_2$	1	1	0
$e_1 - e_3$	1	1	2
$e_2 - e_3$	0	0	0

FIG. 8

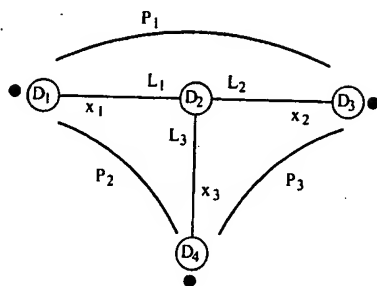


FIG. 9

Flow matrix 1

$$\begin{bmatrix} 1 & 1 & 0 \\ 1 & 0 & 1 \end{bmatrix}$$

Flow matrix 2

$$\begin{bmatrix} 1 & 1 & 0 \\ 1 & 0 & 1 \\ 0 & 1 & 1 \end{bmatrix}$$

FIG. 10

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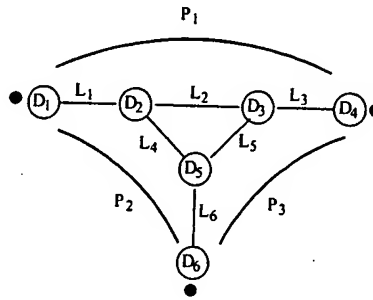


FIG. 11

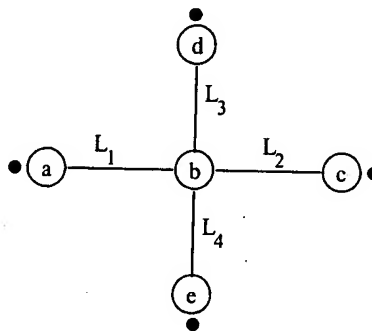


FIG. 12

	L_1	L_2	L_3	L_4
$L_1.L_4$	1	0	0	1
$L_4.L_2$	0	1	0	1
$L_2.L_3$	0	1	1	0
$L_3.L_1$	1	0	1	0

$\{L_1.L_4, L_4.L_2, L_2.L_3, L_3.L_1\}$

	L_1	L_2	L_3	L_4
$L_1.L_2$	1	1	0	0
$L_1.L_3$	1	0	1	0
$L_1.L_4$	1	0	0	1
$L_2.L_3$	0	1	1	0

$\{L_1, L_2, L_3, L_4\}$

FIG. 13

Select_Matrix($G' = (D', I)$): Reduced Network Topology Graph, E : Set of Leaves)

W : Set of worms in G' wrt E , $W \leftarrow \emptyset$

R : Set of paths, $R \leftarrow \emptyset$

Compute P' for G' wrt E

$open \leftarrow P'$

while $open \neq \emptyset$

 select p from $open$

 for each pipe c_i on $p = c_1.c_2 \dots c_{length(p)}$

 if $\exists S \subset open$ such that S makes c_i estimable

 Compute S' which has the original value of each path in S

$R \leftarrow R \cup S'$

$W \leftarrow W \cup \{c_i\}$

 update $open$ and W such that $\forall p' \in open$

p' does not contain any estimable path in W

 else

$c_{i+1} \leftarrow c_i.c_{i+1}$

$open \leftarrow open \setminus \{p\}$

 return W, R

// c_i is removed from paths in $open$

$\leftarrow B1$

FIG. 14

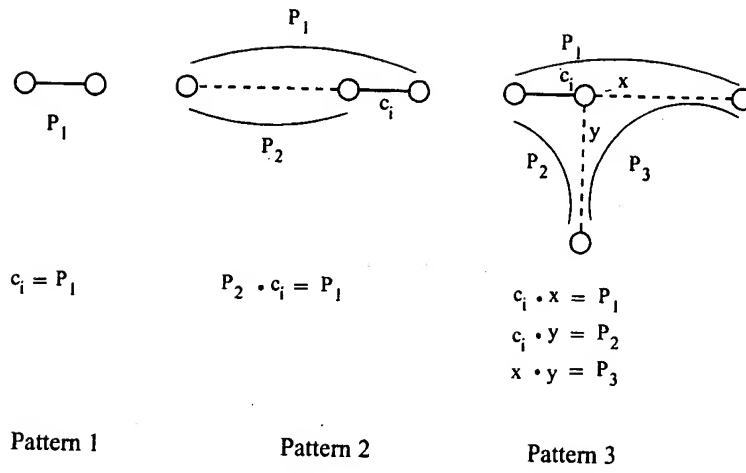


FIG. 15

Compute.EstPaths($G' = (D', I)$): Reduced Network Topology Graph, E : Set of Leaves, P'_{t_i} : End-to-end paths at time t_i)

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 $M$ : A Minimal set of estimable paths for  $G'$  wrt  $E$ ,  $M \leftarrow \emptyset$ 
 $open \leftarrow P'_{t_i}$ 
while  $open \neq \emptyset$ 
    while  $open$  not converged
        select  $p$  from  $open$ 
        for each pipe  $c_i$  on  $p = c_1.c_2 \dots c_{length(p)}$ 
            if  $\exists S \subset open$  such that  $S$  makes  $c_i$  estimable
                 $M \leftarrow M \cup \{c_i\}$ 
                update  $open$  and  $M$  such that  $\forall p' \in open$ 
                     $p'$  does not contain any estimable path in  $M$  and
                     $open \leftarrow open \setminus \{p\}$ 
                    //  $c_i$  is removed from paths in  $open$ 
            else
                abort processing of  $p$ 
        if  $open \neq \emptyset$ 
            select shortest  $p$  in  $open$ 
             $open \leftarrow open \setminus \{p\}$ 
             $M \leftarrow M \cup \{p\}$ 
return  $M$ 

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FIG. 16